

Network Working Group  
Request for Comments: 3630  
Updates: 2370  
Category: Standards Track

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September 2003

## Traffic Engineering (TE) Extensions to OSPF Version 2

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### Abstract

This document describes extensions to the OSPF protocol version 2 to support intra-area Traffic Engineering (TE), using Opaque Link State Advertisements.

### 1. Introduction

This document specifies a method of adding traffic engineering capabilities to OSPF Version 2 [1]. The architecture of traffic engineering is described in [5]. The semantic content of the extensions is essentially identical to the corresponding extensions to IS-IS [6]. It is expected that the traffic engineering extensions to OSPF will continue to mirror those in IS-IS.

The extensions provide a way of describing the traffic engineering topology (including bandwidth and administrative constraints) and distributing this information within a given OSPF area. This topology does not necessarily match the regular routed topology, though this proposal depends on Network LSAs to describe multi-access links. This document purposely does not say how the mechanisms described here can be used for traffic engineering across multiple OSPF areas; that task is left to future documents. Furthermore, no changes have been made to the operation of OSPFv2 flooding; in

particular, if non-TE capable nodes exist in the topology, they MUST flood TE LSAs as any other type 10 (area-local scope) Opaque LSAs (see [3]).

### 1.1. Applicability

Many of the extensions specified in this document are in response to the requirements stated in [5], and thus are referred to as "traffic engineering extensions", and are also commonly associated with MPLS Traffic Engineering. A more accurate (albeit bland) designation is "extended link attributes", as the proposal is to simply add more attributes to links in OSPF advertisements.

The information made available by these extensions can be used to build an extended link state database just as router LSAs are used to build a "regular" link state database; the difference is that the extended link state database (referred to below as the traffic engineering database) has additional link attributes. Uses of the traffic engineering database include:

- o monitoring the extended link attributes;
- o local constraint-based source routing; and
- o global traffic engineering.

For example, an OSPF-speaking device can participate in an OSPF area, build a traffic engineering database, and thereby report on the reservation state of links in that area.

In "local constraint-based source routing", a router R can compute a path from a source node A to a destination node B; typically, A is R itself, and B is specified by a "router address" (see below). This path may be subject to various constraints on the attributes of the links and nodes that the path traverses, e.g., use green links that have unreserved bandwidth of at least 10Mbps. This path could then be used to carry some subset of the traffic from A to B, forming a simple but effective means of traffic engineering. How the subset of traffic is determined, and how the path is instantiated, is beyond the scope of this document; suffice it to say that one means of defining the subset of traffic is "those packets whose IP destinations were learned from B", and one means of instantiating paths is using MPLS tunnels. As an aside, note that constraint-based routing can be NP-hard, or even unsolvable, depending on the nature of the attributes and constraints, and thus many implementations will use heuristics. Consequently, we don't attempt to sketch an algorithm here.

Finally, for "global traffic engineering", a device can build a traffic engineering database, input a traffic matrix and an optimization function, crunch on the information, and thus compute optimal or near-optimal routing for the entire network. The device can subsequently monitor the traffic engineering topology and react to changes by recomputing the optimal routes.

## 1.2. Limitations

As mentioned above, this document specifies extensions and procedures for intra-area distribution of Traffic Engineering information. Methods for inter-area and inter-AS (Autonomous System) distribution are not discussed here.

The extensions specified in this document capture the reservation state of point-to-point links. The reservation state of multi-access links may not be accurately reflected, except in the special case in which there are only two devices in the multi-access subnetwork. Operation over multi-access networks with more than two devices is not specifically prohibited. A more accurate description of the reservation state of multi-access networks is for further study.

This document also does not support unnumbered links. This deficiency will be addressed in future documents; see also [7] and [8].

## 1.3. Conventions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in BCP 14, RFC 2119 [2].

## 2. LSA Format

### 2.1. LSA type

This extension makes use of the Opaque LSA [3].

Three types of Opaque LSAs exist, each of which has a different flooding scope. This proposal uses only Type 10 LSAs, which have an area flooding scope.

One new LSA is defined, the Traffic Engineering LSA. This LSA describes routers, point-to-point links, and connections to multi-access networks (similar to a Router LSA). For traffic engineering purposes, the existing Network LSA is sufficient for describing multi-access links, so no additional LSA is defined for this purpose.





If IS-IS is also active in the domain, this address can also be used to compute the mapping between the OSPF and IS-IS topologies. For example, suppose a router R is advertising both IS-IS and OSPF Traffic Engineering LSAs, and suppose further that some router S is building a single Traffic Engineering Database (TED) based on both IS-IS and OSPF TE information. R may then appear as two separate nodes in S's TED. However, if both the IS-IS and OSPF LSAs generated by R contain the same Router Address, then S can determine that the IS-IS TE LSA and the OSPF TE LSA from R are indeed from a single router.

The router address TLV is type 1, has a length of 4, and a value that is the four octet IP address. It must appear in exactly one Traffic Engineering LSA originated by a router.

#### 2.4.2. Link TLV

The Link TLV describes a single link. It is constructed of a set of sub-TLVs. There are no ordering requirements for the sub-TLVs.

Only one Link TLV shall be carried in each LSA, allowing for fine granularity changes in topology.

The Link TLV is type 2, and the length is variable.

The following sub-TLVs of the Link TLV are defined:

- 1 - Link type (1 octet)
- 2 - Link ID (4 octets)
- 3 - Local interface IP address (4 octets)
- 4 - Remote interface IP address (4 octets)
- 5 - Traffic engineering metric (4 octets)
- 6 - Maximum bandwidth (4 octets)
- 7 - Maximum reservable bandwidth (4 octets)
- 8 - Unreserved bandwidth (32 octets)
- 9 - Administrative group (4 octets)

This memo defines sub-Types 1 through 9. See the IANA Considerations section for allocation of new sub-Types.

The Link Type and Link ID sub-TLVs are mandatory, i.e., must appear exactly once. All other sub-TLVs defined here may occur at most once. These restrictions need not apply to future sub-TLVs. Unrecognized sub-TLVs are ignored.

Various values below use the (32 bit) IEEE Floating Point format. For quick reference, this format is as follows:

Diagram illustrating the IEEE 754 single-precision floating-point format (32 bits):

- Sign (S): 1 bit
- Exponent: 8 bits
- Fraction: 23 bits

The format is divided into four groups of 8 bits each, labeled 0, 1, 2, and 3 at the top. The sign bit is the first bit of group 0. The exponent is the next 8 bits. The fraction is the remaining 23 bits, with the first bit of group 3 being the last bit of the fraction.

S is the sign, Exponent is the exponent base 2 in "excess 127" notation, and Fraction is the mantissa - 1, with an implied binary point in front of it. Thus, the above represents the value:

$$(-1)^{(S)} \cdot 2^{(Exponent-127)} \cdot (1 + Fraction)$$

For more details, refer to [4].

## 2.5. Sub-TLV Details

### 2.5.1. Link Type

The Link Type sub-TLV defines the type of the link:

- 1 - Point-to-point
- 2 - Multi-access

The Link Type sub-TLV is TLV type 1, and is one octet in length.

### 2.5.2. Link ID

The Link ID sub-TLV identifies the other end of the link. For point-to-point links, this is the Router ID of the neighbor. For multi-access links, this is the interface address of the designated router. The Link ID is identical to the contents of the Link ID field in the Router LSA for these link types.

The Link ID sub-TLV is TLV type 2, and is four octets in length.

### 2.5.3. Local Interface IP Address

The Local Interface IP Address sub-TLV specifies the IP address(es) of the interface corresponding to this link. If there are multiple local addresses on the link, they are all listed in this sub-TLV.

The Local Interface IP Address sub-TLV is TLV type 3, and is 4N octets in length, where N is the number of local addresses.

#### 2.5.4. Remote Interface IP Address

The Remote Interface IP Address sub-TLV specifies the IP address(es) of the neighbor's interface corresponding to this link. This and the local address are used to discern multiple parallel links between systems. If the Link Type of the link is Multi-access, the Remote Interface IP Address is set to 0.0.0.0; alternatively, an implementation MAY choose not to send this sub-TLV.

The Remote Interface IP Address sub-TLV is TLV type 4, and is 4N octets in length, where N is the number of neighbor addresses.

#### 2.5.5. Traffic Engineering Metric

The Traffic Engineering Metric sub-TLV specifies the link metric for traffic engineering purposes. This metric may be different than the standard OSPF link metric. Typically, this metric is assigned by a network administrator.

The Traffic Engineering Metric sub-TLV is TLV type 5, and is four octets in length.

#### 2.5.6. Maximum Bandwidth

The Maximum Bandwidth sub-TLV specifies the maximum bandwidth that can be used on this link, in this direction (from the system originating the LSA to its neighbor), in IEEE floating point format. This is the true link capacity. The units are bytes per second.

The Maximum Bandwidth sub-TLV is TLV type 6, and is four octets in length.

#### 2.5.7. Maximum Reservable Bandwidth

The Maximum Reservable Bandwidth sub-TLV specifies the maximum bandwidth that may be reserved on this link, in this direction, in IEEE floating point format. Note that this may be greater than the maximum bandwidth (in which case the link may be oversubscribed). This SHOULD be user-configurable; the default value should be the Maximum Bandwidth. The units are bytes per second.

The Maximum Reservable Bandwidth sub-TLV is TLV type 7, and is four octets in length.



### 2.5.8. Unreserved Bandwidth

The Unreserved Bandwidth sub-TLV specifies the amount of bandwidth not yet reserved at each of the eight priority levels in IEEE floating point format. The values correspond to the bandwidth that can be reserved with a setup priority of 0 through 7, arranged in increasing order with priority 0 occurring at the start of the sub-TLV, and priority 7 at the end of the sub-TLV. The initial values (before any bandwidth is reserved) are all set to the Maximum Reservable Bandwidth. Each value will be less than or equal to the Maximum Reservable Bandwidth. The units are bytes per second.

The Unreserved Bandwidth sub-TLV is TLV type 8, and is 32 octets in length.

### 2.5.9. Administrative Group

The Administrative Group sub-TLV contains a 4-octet bit mask assigned by the network administrator. Each set bit corresponds to one administrative group assigned to the interface. A link may belong to multiple groups.

By convention, the least significant bit is referred to as 'group 0', and the most significant bit is referred to as 'group 31'.

The Administrative Group is also called Resource Class/Color [5].

The Administrative Group sub-TLV is TLV type 9, and is four octets in length.

## 3. Elements of Procedure

Routers shall originate Traffic Engineering LSAs whenever the LSA contents change, and whenever otherwise required by OSPF (an LSA refresh, for example). Note that this does not mean that every change must be flooded immediately; an implementation MAY set thresholds (for example, a bandwidth change threshold) that trigger immediate flooding, and initiate flooding of other changes after a short time interval. In any case, the origination of Traffic Engineering LSAs SHOULD be rate-limited to at most one every MinLSInterval [1].

Upon receipt of a changed Traffic Engineering LSA or Network LSA (since these are used in traffic engineering calculations), the router should update its traffic engineering database. No Shortest Path First (SPF) or other route calculations are necessary.

#### 4. Compatibility Issues

There should be no interoperability issues with routers that do not implement these extensions, as the Opaque LSAs will be silently ignored.

The result of having routers that do not implement these extensions is that the traffic engineering topology will be missing pieces. However, if the topology is connected, TE paths can still be calculated and ought to work.

#### 5. Security Considerations

This document specifies the contents of Opaque LSAs in OSPFv2. As Opaque LSAs are not used for SPF computation or normal routing, the extensions specified here have no affect on IP routing. However, tampering with TE LSAs may have an effect on traffic engineering computations, and it is suggested that any mechanisms used for securing the transmission of normal OSPF LSAs be applied equally to all Opaque LSAs, including the TE LSAs specified here.

Note that the mechanisms in [1] and [9] apply to Opaque LSAs. It is suggested that any future mechanisms proposed to secure/authenticate OSPFv2 LSA exchanges be made general enough to be used with Opaque LSAs.

#### 6. IANA Considerations

The top level Types in a TE LSA, as well as Types for sub-TLVs for each top level Type, have been registered with IANA, except as noted.

Here are the guidelines (using terms defined in [10]) for the assignment of top level Types in TE LSAs:

- o Types in the range 3-32767 are to be assigned via Standards Action.
- o Types in the range 32768-32777 are for experimental use; these will not be registered with IANA, and MUST NOT be mentioned by RFCs.
- o Types in the range 32778-65535 are not to be assigned at this time. Before any assignments can be made in this range, there MUST be a Standards Track RFC that specifies IANA Considerations that covers the range being assigned.

The guidelines for the assignment of types for sub-TLVs in a TE LSA are as follows:

- o Types in the range 10-32767 are to be assigned via Standards Action.
- o Types in the range 32768-32777 are for experimental use; these will not be registered with IANA, and MUST NOT be mentioned by RFCs.
- o Types in the range 32778-65535 are not to be assigned at this time. Before any assignments can be made in this range, there MUST be a Standards Track RFC that specifies IANA Considerations that covers the range being assigned.

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## 8. References

### 8.1. Normative References

- [1] Moy, J., "OSPF Version 2", STD 54, RFC 2328, April 1998.
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### 8.2. Informative References

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## Acknowledgement

Funding for the RFC Editor function is currently provided by the Internet Society.

